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Random Spanning Trees and the Matrix Tree Theorem.

We are now going to prove some amazing connections between eigenvalues and spanning trees.

We start with some determinant facts.

Leibniz formula. For an $n \times n$ matrix A ,

$$\det A = \sum_{\pi \in S_n} (\text{sgn}(\pi) \prod_{i=1}^n A_{i,\pi(i)})$$

where S_n is the group of permutations of n letters.

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Definition. The sign of a permutation is defined by $\text{sgn}(\pi) = (-1)^{\# \text{of } 2\text{-cycles}}$.

Fact. Although the decomposition of π into 2-cycles (or transpositions) is not unique, the parity of the # of 2-cycles is the same in all decomposition of π .

Product formula: $\det(AB) = \det A \det B$.

Invariance under row operations.

For every c , if $A = (\overset{\uparrow}{\vec{a}_1} \cdots \overset{\uparrow}{\vec{a}_n})$, then

$$\det(\overset{\uparrow}{\vec{a}_1} \cdots \overset{\uparrow}{\vec{a}_n}) = \det(\overset{\uparrow}{\vec{a}_1} \cdots \overset{\uparrow}{\vec{a}_{n-1}} \overset{\uparrow}{\vec{a}_n + c\vec{a}_1})$$

We will also use the fact that $\det(A)$ is linear in every column of A .

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Determinant as volume. If $P(A)$
 is the polytope whose corners are
 $\left\{ \sum_{i \in S} \vec{a}_i \mid S \subseteq \{1, \dots, n\} \right\}$ (the parallelepiped
 with axes $\vec{a}_1, \dots, \vec{a}_n$.

Now let $\Pi_{\vec{a}_1}$ be the projection

$$\Pi_{\vec{a}_1}(\vec{v}) = \vec{v} - \frac{\langle \vec{v}, \vec{a}_1 \rangle}{\langle \vec{a}_1, \vec{a}_1 \rangle} \vec{a}_1$$

onto the ~~perp~~ hyperspace normal to \vec{a}_1 .
 We know that

$$\begin{aligned} \det(\vec{a}_1, \dots, \vec{a}_n) &= \det(\vec{a}_1, \Pi_{\vec{a}_1} \vec{a}_2, \dots, \Pi_{\vec{a}_1} \vec{a}_n) \\ &= \|\vec{a}_1\| \cdot \text{volume of } P(\alpha_2 \Pi_{\vec{a}_1} \vec{a}_2, \dots, \Pi_{\vec{a}_1} \vec{a}_n) \end{aligned}$$

but what is this volume?

④ We know that

$$\det(xI - A) = \sum_{k=0}^n x^{n-k} (-1)^k \sigma_k(A)$$

where $\sigma_k(A)$ is the k -th elementary symmetric polynomial in the eigenvalues of A .

$$\sigma_k(A) := \sum_{\substack{S = \{1, \dots, n\} \\ |S|=k}} \prod_{i \in S} \lambda_i$$

$$\sigma_1(A) = \lambda_1 + \dots + \lambda_n = \text{tr } A$$

$$\sigma_0(A) = 1 \quad (\text{by convention})$$

:

$$\sigma_n(A) = \lambda_1 \cdots \lambda_n = \det A.$$

Notice that if we change A to a similar matrix XAX^{-1} , this doesn't change the eigenvalues. (old homework)

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so it doesn't change $\sigma_1, \dots, \sigma_n$ either.

We also proved that

$$\sigma_K(A) = \sum_{|S|=K} \det(A(S,S))$$

= sum of determinants of
principal $K \times K$ minors.

With all this in hand,

$$\text{vol}(\Pi_{\vec{a}_1} \vec{a}_2, \dots, \Pi_{\vec{a}_{n-1}} \vec{a}_n) = \sigma_{n-1}(\vec{0}, \Pi_{\vec{a}_1} \vec{a}_2, \dots, \Pi_{\vec{a}_{n-1}} \vec{a}_n)$$

because $\sigma_{n-1}(\vec{0}, \Pi_{\vec{a}_1} \vec{a}_2, \dots, \Pi_{\vec{a}_{n-1}} \vec{a}_n)$ has
only one nonzero term, given by
the product of the $n-1$ nonzero
eigenvalues.

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We can also compute

$$\text{vol} \left(\Pi_{\vec{a}_1} \vec{a}_2, \dots, \Pi_{\vec{a}_1} \vec{a}_n \right)^2$$

$$= \det \left((\Pi_{\vec{a}_1} \vec{a}_2 \dots \Pi_{\vec{a}_1} \vec{a}_n)^T (\Pi_{\vec{a}_1} \vec{a}_2 \dots \Pi_{\vec{a}_1} \vec{a}_n) \right)$$

where the matrix inside the determinant is the Gramian matrix of dot products of these $n-1$ vectors and the det is called the Gramian determinant.

As a last fact, if A and B are symmetric matrices of rank K with the same nullspace,

$$\sigma_K(AB) = \sigma_K(A)\sigma_K(B)$$

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Matrix Tree Theorem.

Let $G = (V, E, \omega)$ be a ^{connected} [^]weighted graph,

$$\sigma_{n-1}(L_G) = n \sum_{\substack{\text{Spanning} \\ \text{trees of } G}} \prod_{\substack{a \sim b \\ \in T}} \omega(a \sim b)$$

Proof. Suppose that G is a tree.

For each vertex $a \in G$, let $S_a = V - \{a\}$.

We Know

$$\sigma_{n-1}(L_G) = \sum_{a \in V} \det(L_G(S_a, S_a))$$

Now we can always write

$$L_G = U^T W U$$

\uparrow \uparrow
 $V \times E$ $E \times V$

adjacency matrix

~~boundary matrix $B \in \mathbb{R}^{E \times V}$~~

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where W is the $E \times E$ diagonal matrix of edge weights.

If we recall that also

$$L_G = \underbrace{(U^T W^{1/2})}_{B^*} \underbrace{(W^{1/2} U)}_{B^T}$$

\uparrow boundary
 $V \times E$
 transpose

\uparrow boundary matrix
 $E \times V$

We have

~~det~~

$$L_G(S_a, S_a) = \underbrace{B(S_a, :)}_{\text{a square matrix}} B^T(:, S_a)$$

b/c a tree has $V+1$ edges (always)

so

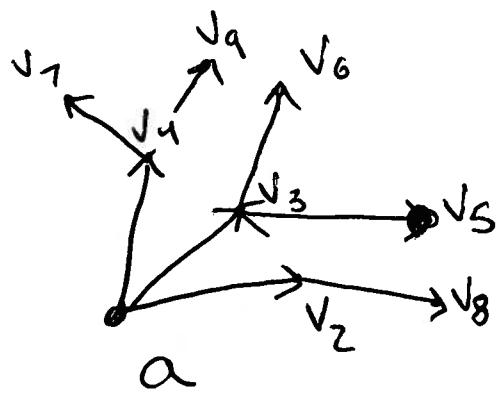
$$\det(L_G(S_a, S_a)) = \det(B(S_a, :))^2$$

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Now we have already shown that renumbering vertices and edges permutes row and columns of B .

This transformation can't change $\det B(S_a, \text{all})$.

So let's number a particular way.



Start at $a = v_1$, and number vertices by distance from a (breaking ties randomly). ~~Orient each edge by distance from a so that the head~~

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Now each v_i has one edge $v_j \rightarrow v_i$ with $j < i$. We call this edge e_{i-1} .

Now we know

$$B_{ij} = \begin{cases} +\sqrt{\omega_j} & \text{if vertex } i \text{ is the head of edge } j \\ -\sqrt{\omega_j} & \text{if vertex } i \text{ is the tail of edge } j \\ 0 & \text{otherwise} \end{cases}$$

so in our new numbering

$$B_{i,i-1} = \pm \sqrt{\omega_{i-1}}, \text{ because } v_i \text{ is either head or tail of } e_{i-1}$$

$$B_{i,k} = 0 \text{ if } k > i-1, \text{ because the other vertex } v_j \text{ incident edge } e_{i-1} \text{ has } j < i$$

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If we now take

$$B(S_a; \text{all}) = B(\{2, \dots, n\}, \text{all})$$

we see that

$$B(\{2, \dots, n\}, \text{all})_{i,i} = B_{i+1,i} = \pm \sqrt{\omega_i}$$

and

$$B(\{2, \dots, n\}, \text{all})_{i,k} = B_{i+1,k} = 0 \text{ if } k > i$$

so

$B(\{2, \dots, n\}, \text{all})$ is lower-triangular
and has determinant $\pm \prod_i \sqrt{\omega_i}$.

Thus

$$\det(B(S_a; \text{all}))^2 = \prod_i \omega_i.$$

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This proves the theorem if G is a tree.

Now in general, let

$$L_G = BB^T, \text{ where } B = U^T W^{1/2}$$

We Know

\downarrow ~~verts × verts~~

$$\begin{aligned} \sigma_{n-1}(L_G) &= \sigma_{n-1}(B^T B) \quad \begin{array}{l} \text{edges} \times \text{edges} \\ (\cancel{\text{verts}} \times \cancel{\text{verts}}) \\ \cancel{\text{verts}}, \text{ because} \end{array} \\ &= \sigma_{n-1}(B^T B) \quad \begin{array}{l} \text{the eigenvalues} \\ \text{are the same} \\ (\text{up to some zeros}) \end{array} \end{aligned}$$

$$\begin{aligned} &= \sum_{\substack{|S|=\text{verts}-1 \\ S \subseteq E}} \sigma_{n-1}(\underbrace{B^T(S; \text{all}) B(\text{all}; S)}_{\text{principal minor}}) \\ &\qquad \qquad \qquad B^T B(S, S) \end{aligned}$$

$$\begin{aligned} &= \sum_{\substack{|S|=\text{verts}-1 \\ S \subseteq E}} \sigma_{n-1}(B(\text{all}, S) B^T(S, \text{all})) \\ &\qquad \qquad \qquad \begin{array}{l} \uparrow \\ \text{again eigenvalues are} \\ \text{same up to zeros!} \end{array} \end{aligned}$$

But we've done something very clever here!

$B(\text{all}, S)$ = $\text{verts} \times |S|$ matrix

where $B_{ij} = \begin{cases} +\sqrt{\omega_j} & \text{if } v_i \text{ is head} \\ & \text{of } e_j \\ -\sqrt{\omega_j} & \text{if } v_i \text{ is tail} \\ & \text{of } e_j \\ 0, & \text{otherwise} \end{cases}$

δ is the boundary operator for the subgraph of G containing only edges in S . ~~This~~ Call this G_S . Then

$$B(\text{all}; S) B^T(S; \text{all}) = L_{G_S}$$

so we have shown

$$\sigma_{n-1}(L_G) = \sum_{|S|=\text{verts}-1} \sigma_{n-1}(L_{G_S})$$

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But if G_s is connected, it is a tree, and by our previous argument, $\sigma_{n-1}(L_{G_s}) = \sum_{e_j \in S} w_j$.

If G_s is disconnected, it has > 1 zero eigenvalue and $\sigma_{n-1}(L_{G_s}) = 0$.

Thus

$$\sigma_{n-1}(L_G) = \sum_{\substack{T \text{ is a spanning tree of } G \\ e_j \in T}} \prod_{e_i \in T} w_j \quad \square.$$