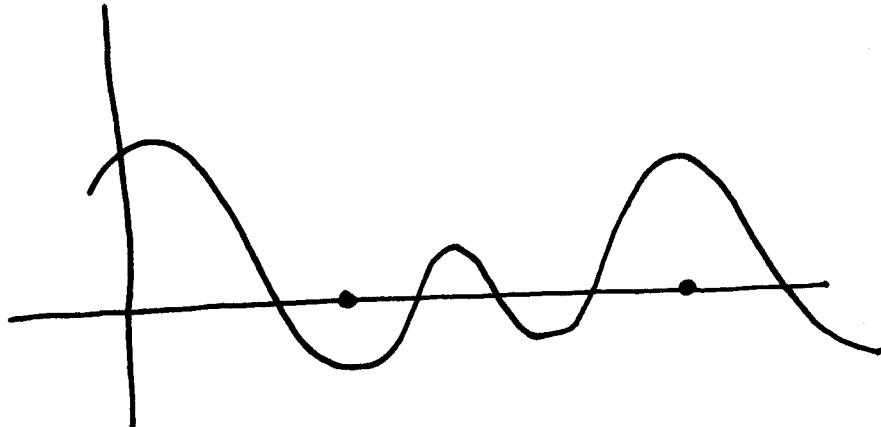


3.1. Finding Roots of Functions.

①

We are interested in solving equations of the form $f(x)=0$ for x .



Idea 1. Given an interval $[a,b]$ where $\text{sign } f(a) \neq \text{sign } f(b)$, we compute $f(c)$ where $c = \frac{a+b}{2}$ and restrict our attention to either $[a,c]$ or $[c,b]$ depending on $\text{sign } f(c)$.

This is called the bisection method. *

<Mathematica demonstration>

(2)

After n steps, the error in the computed position of the root is at most $\frac{b-a}{2^{n+1}}$.

Definition. If $\{x_n\} \rightarrow x$, then the sequence has linear convergence if $\exists C \in [0, 1)$ so that

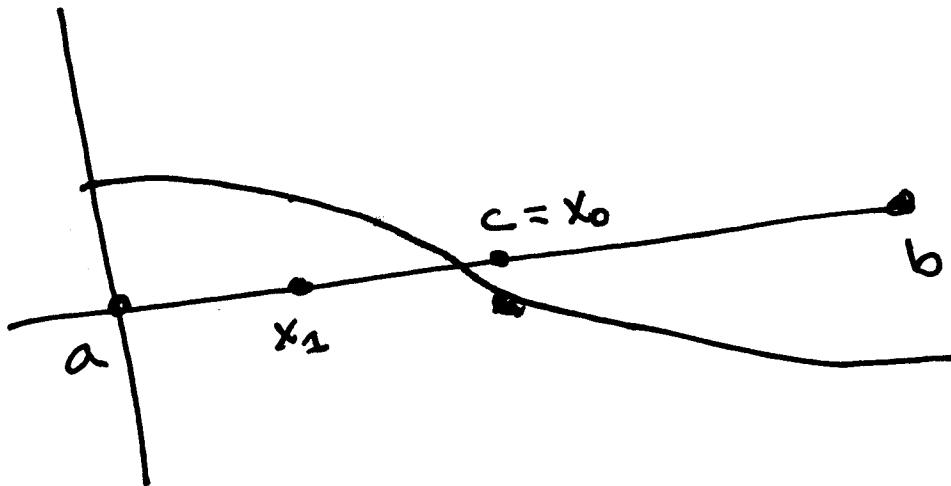
$$|x_{n+1} - x| \leq C|x_n - x|$$

Lemma. If $\{x_n\} \rightarrow x$ linearly with constant C , then $|x_{n+1} - x| \leq AC^n$, where $A = \frac{\max|x_1 - x|}{|x_1 - x|}$.

Question. Does the bisection method converge linearly? (We take the sequence to be the sequence of ~~end~~ ^{mid} points, and x to be whatever root the method converges to.)

Consider

(3)



On step 1, we replace our initial guess of $a+b/2$ by $\frac{a+\frac{a+b}{2}}{2} = \frac{3}{4}a + \frac{1}{4}b = x_1$.

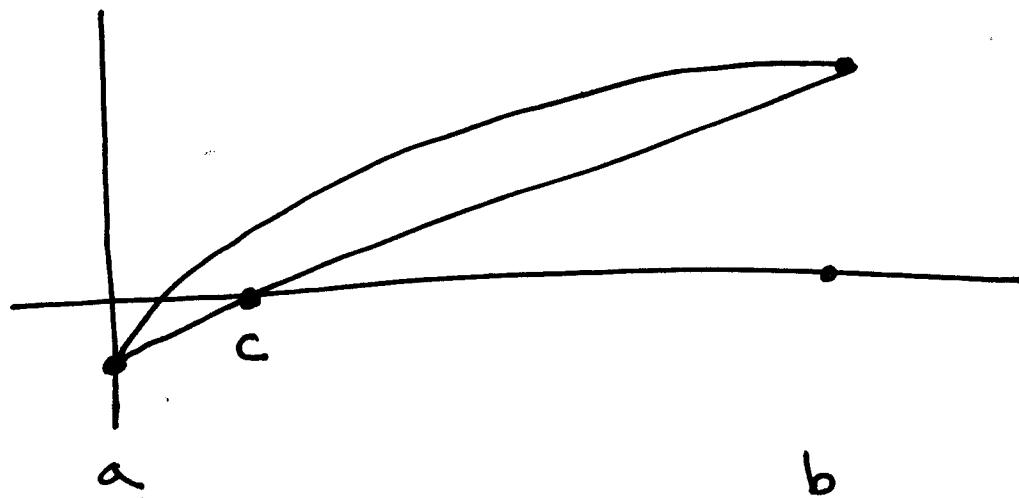
But x_0 was actually closer to the root than x_1 !

We see that bisection is not guaranteed to improve at each step.

On the other hand, bisection does give the conclusion of the Lemma, which is almost equally useful in practice.

(4)

We can improve the bisection method by changing our choice of midpoint.

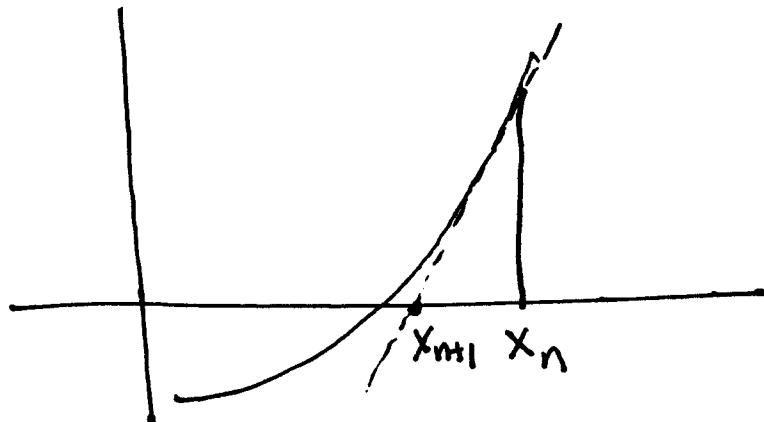


The "false position" method guesses that the function is well approximated by the secant line through $(a, f(a))$ and $(b, f(b))$ and chooses the guess for the new zero accordingly.

This can have linear (and even superlinear) convergence if the details are handled right... we will return to this method soon!

5.

What if we can calculate a derivative of our function?



Estimating where the ~~tangent~~ tangent line crosses the x-axis is the basis for Newton's Method.

Doing the algebra establishes that

$$x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)}$$

< newton-method.n b demonstration >

(6)

We saw that the number of correct digits increases exponentially. To be more precise, we will prove

Definition. We say $\{x_n\} \rightarrow x$ quadratically if $|x_{n+1} - x| \leq c|x_n - x|^2$ for some c .

Observe that if x_n has K correct decimal digits, then $|x_n - x| < 10^{-K}$, so

$|x_{n+1} - x| < c(10^{-K})^2 = c10^{-2K}$, and the number of correct digits has ~~almost~~ approximately doubled (depending on c).

Newton's Method Theorem. If f, f', f'' are continuous in a neighborhood of a root r of f , and $f'(r) \neq 0$, there is a neighborhood N_δ of r of radius δ so that if $x_0 \in N_\delta$ then all $x_n \in N_\delta$ and

$$|r - x_{n+1}| \leq c(\delta) |r - x_n|^2$$

for some c depending on f and δ (given below).

(7)

Proof. Let $e_n = r - x_n$. We know

$$\begin{aligned}
 e_{n+1} &= r - x_{n+1} = r - \left(x_n - \frac{f(x_n)}{f'(x_n)} \right) \\
 &= (r - x_n) + \frac{f(x_n)}{f'(x_n)} \\
 &= e_n + \frac{f(x_n)}{f'(x_n)} \\
 &= \frac{e_n f'(x_n) + f(x_n)}{f'(x_n)}.
 \end{aligned}$$

Now let's Taylor expand f around x_n .
We know

$$0 = f(r) = f(x_n + e_n) = f(x_n) + e_n f'(x_n) + \frac{1}{2} e_n^2 f''(\xi_n)$$

where $\xi_n \in [x_n, r]$. This means

$$e_{n+1} = -\frac{1}{2} \cancel{\frac{f''(\xi_n)}{f'(x_n)}} e_n^2,$$

which is almost what we want.

Observe that we can define a function

$$c(\delta) = \frac{1}{2} \frac{\max_{N\delta} |f''(x)|}{\min_{N\delta} |f'(x)|}$$

which is finite for small enough δ . In fact, we can choose δ small enough that

$$\delta c(\delta) < 1$$

since as $\delta \rightarrow 0$, $c(\delta) \rightarrow \frac{f''(r)}{f'(r)}$. Now all we

have to observe is that if $x_n \in N\delta$,

$$\left| \frac{1}{2} \frac{f''(\xi_n)}{f'(x_n)} \right| \leq c(\delta)$$

Since, x_n, ξ_n are in $N\delta$. In this case

$$|e_{n+1}| \leq c(\delta) e_n^2 \leq \delta c(\delta) e_n < e_n,$$

so x_{n+1} is in $N\delta$ as well. So if we choose x_0 in $N\delta$, all subsequent x_n are also in $N\delta$.

We now only have to show that $\{x_n\} \rightarrow r$. (9)

Observe

$$|e_n| < \delta c(\delta) |e_{n-1}| < \dots < (\delta c(\delta))^{n+1} e_0.$$

Since $\delta c(\delta) < 1$, this means $\{e_n\} \rightarrow 0$, as desired.

A very cool extension of Newton's method is commonly used to solve systems of nonlinear equations.

Idea: Given the system

$$f_1(x_1, \dots, x_N) = 0$$

⋮

$$f_N(x_1, \dots, x_N) = 0$$

we can think of this as a map

$$F: \mathbb{R}^n \rightarrow \mathbb{R}^n, \text{ given by } F(x_1, \dots, x_n) = (f_1, \dots, f_n).$$

(10)

We then write down Newton's method,
recalling that ~~the~~ =

$$F(\vec{x} + \vec{h}) \approx F(\vec{x}) + DF_{\vec{x}}(\vec{h}) = L(\vec{h})$$

where

$$DF_{\vec{x}} = \begin{bmatrix} \frac{\partial f_1}{\partial x_1} & \frac{\partial f_1}{\partial x_2} & \cdots & \frac{\partial f_1}{\partial x_n} \\ \vdots & \vdots & & \vdots \\ \frac{\partial f_n}{\partial x_1} & \cdots & \cdots & \frac{\partial f_n}{\partial x_n} \end{bmatrix}$$

Now the linear approximation $L(\vec{h})$ vanishes when \vec{h} solves the linear system

$$(DF_{\vec{x}}) \vec{h} = -F(\vec{x}).$$

So we let \vec{h}_n be the solution to

$$(DF_{\vec{x}_n}) \vec{h} = -F(\vec{x}_n)$$

and set

$$\begin{aligned}\vec{x}_{n+1} &= \vec{x}_n + \vec{h}_n \\ &= \vec{x}_n - (\text{DF}_x)^{-1} F(\vec{x}_n).\end{aligned}$$

Notice that this only works when DF_x is a nonsingular matrix, just as 1-d Newton's method works only for $f' \neq 0$.

Example. <computer>

One of the interesting features of Newton's method is how it chooses a final point to converge to.

<demonstration with solving $z^3=1$ in complex plane>