

Conjugate Directions.

(1)

We've seen that neither coordinate directions* nor gradient descent are very effective when approaching a very "elliptical" minimum.

What's the problem? Observe

Lemma. If we minimize f along direction \vec{d} from \vec{x} at the minimum, $\nabla f \perp \vec{d}$.

Proof. $\langle \nabla f, \vec{d} \rangle$ is the derivative of f along the line. \square

Now we note that around any P ,

$$f(\vec{p} + \vec{x}) = f(\vec{p}) + \nabla f(\vec{p}) \cdot \vec{x} + \frac{1}{2} \vec{x}^T H(\vec{p}) \vec{x} + \dots$$

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* Note that if your function is a sum or (more subtly) a product of functions of the individual variables x_1, \dots, x_n , then coordinate search works beautifully, solving the problem exactly in n steps.

So if $G(\vec{p}) = -b$, and $H(\vec{p}) = A$,
we can compute

$$\nabla f(\vec{p} + \vec{x}) = -b + A\vec{x}.$$

Now we observe that this implies
we can find the \vec{x} so that $\nabla f(\vec{p} + \vec{x}) = 0$
by solving

$$A\vec{x} - \vec{b} = 0.$$

(and there are methods which do this!).

Now how does ∇f change as we
move in direction \vec{v} ? Well, we expect

$$\begin{aligned} \Delta(\nabla f) &\approx \Delta(-b + A\vec{x}) \\ &= A\vec{v} \end{aligned}$$

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③

Now when we minimized along \vec{u} , we arrived at a point where $\langle \nabla f, \vec{u} \rangle = 0$. In our search for a point where $\nabla f = 0$, we'd like to keep the property that $\langle \nabla f, \vec{u} \rangle = 0$ even when we move along \vec{v} . For that to work, we need the change ~~is~~

$$D_{\vec{v}} (\langle \nabla f, \vec{u} \rangle) = 0$$

or

$$\langle \vec{u}, D_{\vec{v}} (\nabla f) \rangle = 0$$

or (approximately)

$$\langle \vec{u}, A \vec{v} \rangle = 0$$

④

Definition. We say that \vec{u} and \vec{v} are conjugate (with respect to A) if $\langle \vec{u}, A\vec{v} \rangle = 0$.

Note: If A is positive-definite (as it is at a ~~min~~ nondegenerate minimum of $f!$), we note that this is just the condition that \vec{u} and \vec{v} are orthogonal in the metric generated by A .

Meta-algorithm. Generate ~~a~~ a set of n (linearly independent) conjugate directions and minimize along each in turn.

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Fletcher, Theorem 2.4.1.
Proposition. If f is quadratic and the Hessian of f is positive definite, the meta-algorithm finds the exact min (in n steps).

Proof. In general, \exists some \vec{x}' and c' so that any quadratic function of \vec{x} can be written

$$q(\vec{x}) = \frac{1}{2} (\vec{x} - \vec{x}')^T A (\vec{x} - \vec{x}') + c'$$

It's not hard to see that

1) A is the Hessian of q

2) \vec{x}' is the ~~min~~ ^{imizer} of q if A is pos. def.

Suppose $\vec{s}_1, \dots, \vec{s}_n$ are a set of conjugate directions. We first show that the \vec{s}_i form a basis for \mathbb{R}^n .

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Suppose $(\omega \log)$ not. Then

$$\vec{s}_1 = \sum_{i=2}^n a_i \vec{s}_i,$$

~~so~~ where $a_2 \neq 0$ (we could always reorder to get this).

$$\vec{s}_1^T A \vec{s}_2 = \left(\sum a_i \vec{s}_i \right)^T A \vec{s}_2$$

$$= \underbrace{(a_2 \vec{s}_2)^T A \vec{s}_2}_{\neq 0, \text{ since } a_2 \neq 0, A \text{ is pos. def.}} + \underbrace{\sum_{i=3}^n (a_i \vec{s}_i)^T A \vec{s}_2}_{0, \text{ since the } \vec{s}_i \text{ are conjugate to } \vec{s}_2}.$$

$\neq 0$, since $a_2 \neq 0$, A is pos. def.

0 , since the ~~the~~ \vec{s}_i are conjugate to \vec{s}_2

On the other hand,

$$\vec{s}_1^T A \vec{s}_2 = 0, \text{ since } s_1 \text{ is conjugate to } s_2.$$

This means ~~we can write~~ that if we start the meta-algorithm at $\vec{x}^{(1)}$ then we can write the minimizer as

$$\vec{x}' = \vec{x}^{(1)} + \sum a_i' \vec{s}_i$$

and any other point

$$\vec{X} = \vec{X}^{(1)} + \sum a_i \vec{s}_i.$$

Thus we can write q in terms of the vector $\vec{a} = (a_1, \dots, a_n)$ as

$$q(\vec{a}) = \frac{1}{2} (\vec{a} - \vec{a}') S^T A S (\vec{a} - \vec{a}') + c!$$

where $S = \begin{pmatrix} \uparrow \vec{s}_1 & \uparrow \vec{s}_2 & \dots & \uparrow \vec{s}_n \\ \downarrow & \downarrow & & \downarrow \end{pmatrix}$. Now the

\vec{s}_i are conjugate, so

$S^T A S = D$ ← a diagonal matrix with diagonal entries

$$d_i = \vec{s}_i^T A \vec{s}_i.$$

so

$$q(\vec{a}) = \frac{1}{2} \sum (a_i - a'_i)^2 d_i$$

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Now line search along the ~~\vec{a}_i~~ \vec{s}_i directions amounts to coordinate search on the a_i coordinates, which is guaranteed to converge, since $q(\vec{a})$ is the sum of functions of the individual a_i . \square

Clever! But how can we take advantage of this method? If we knew the Hessian, we could generate conjugate directions by Gram-Schmidt relative to the metric generated by the Hessian. But usually, we don't know the Hessian, so we need a little more cleverness.

This will lead us directly to the "conjugate-gradient" methods, which are based on the following theorem.

⑨

Orthogonalization Theorem.

Suppose A is a symmetric positive definite matrix. Let \vec{g}_0 be any vector and $\vec{h}_0 = \vec{g}_0$. Define a sequence of vectors \vec{h}_i, \vec{g}_i by

$$\vec{g}_{i+1} = \vec{g}_i - \lambda_i A \vec{h}_i$$

$$\vec{h}_{i+1} = \vec{g}_{i+1} - \gamma_i \vec{h}_i$$

where we choose λ_i and γ_i so that $\vec{g}_{i+1} \cdot \vec{g}_i = 0$, $\vec{h}_{i+1}^T A \vec{h}_i = 0$. Then for all $i \leq n$, $\{\vec{g}_0, \dots, \vec{g}_i\}$ is a set of mutually orthogonal vectors and

$\{\vec{h}_0, \dots, \vec{h}_i\}$ is a set of mutally conjugate vectors.

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We note that it's easy to see

$$\lambda_i = \frac{\vec{g}_i \cdot \vec{g}_i}{\vec{g}_i^T A \vec{h}_i}, \quad \gamma_i = \frac{\vec{g}_{i+1}^T A \vec{h}_i}{\vec{h}_i^T A \vec{h}_i}.$$

Check:

$$\begin{aligned} \vec{g}_{i+1} \cdot \vec{g}_i &= \left(\vec{g}_i - \frac{\vec{g}_i \cdot \vec{g}_i}{\vec{g}_i^T A \vec{h}_i} A \vec{h}_i \right) \cdot \vec{g}_i \\ &= \cancel{\vec{g}_i \cdot \vec{g}_i} - \frac{\vec{g}_i \cdot \vec{g}_i}{\vec{g}_i^T A \vec{h}_i} \vec{g}_i^T A \vec{h}_i = 0. \end{aligned}$$

$$\begin{aligned} \vec{h}_{i+1}^T A \vec{h}_i &= \left(\vec{g}_{i+1} + \left(\frac{\vec{g}_{i+1}^T A \vec{h}_i}{\vec{h}_i^T A \vec{h}_i} \right) \vec{h}_i \right)^T A \vec{h}_i \\ &= \vec{g}_{i+1}^T A \vec{h}_i - \vec{g}_{i+1}^T A \vec{h}_i \\ &= 0. \end{aligned}$$

Corollary. We claim

$$\gamma_i = -\frac{\vec{g}_{i+1} \cdot \vec{g}_{i+1}}{\vec{g}_i \cdot \vec{g}_i} = -\frac{(\vec{g}_{i+1} - \vec{g}_i) \cdot \vec{g}_{i+1}}{\vec{g}_i \cdot \vec{g}_i}$$

and

$$\lambda_i = \frac{\vec{g}_i \cdot \vec{h}_i}{\vec{h}_i^T A \vec{h}_i}.$$

Proof. Recall that

$$\vec{g}_{i+1} = \vec{g}_i - \lambda_i A \vec{h}_i$$

so $\frac{\vec{g}_{i+1} - \vec{g}_i}{-\lambda_i} = A \vec{h}_i$. Thus

$$\gamma_i = \frac{\vec{g}_{i+1} (A \vec{h}_i)}{\vec{h}_i^T A \vec{h}_i} = \frac{\vec{g}_{i+1} \cdot (\vec{g}_{i+1} - \vec{g}_i)}{-\lambda_i \vec{h}_i^T A \vec{h}_i}$$

Now

$$-\lambda_i \vec{h}_i^T A \vec{h}_i = \frac{-\vec{g}_i \cdot \vec{g}_i}{\vec{g}_i^T A \vec{h}_i} \vec{h}_i^T A \vec{h}_i$$

We now consider

$$\frac{h_i^T A h_i}{g_i^T A h_i}$$

Now

$$\vec{g}_i = \vec{h}_i - \gamma_{i-1} \vec{h}_{i-1}$$

so by conjugacy of the h_i , we have

$$\vec{g}_i^T A \vec{h}_i = \vec{h}_i^T A \vec{h}_i.$$

Thus we get

$$\gamma_i = \frac{\vec{g}_{i+1} \cdot (\vec{g}_{i+1} - \vec{g}_i)}{-\vec{g}_i \cdot \vec{g}_i} = - \frac{\vec{g}_{i+1} \cdot \vec{g}_{i+1}}{\vec{g}_i \cdot \vec{g}_i}.$$

We now tackle

$$\lambda_i = \frac{\vec{g}_i \cdot \vec{g}_i}{\vec{g}_i^T A \vec{h}_i} = \frac{\vec{g}_i \cdot \vec{g}_i}{\vec{g}_i \cdot (\vec{g}_{i+1} - \vec{g}_i)}$$

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Now we see

$$\begin{aligned}\vec{g}_i^T (A \vec{h}_i) &= (\vec{h}_i - \gamma_{i-1} \vec{h}_{i-1})^T A \vec{h}_i \\ &= \vec{h}_i^T A \vec{h}_i\end{aligned}$$

and
~~so~~ we have

$$\begin{aligned}\vec{g}_i \cdot \vec{g}_i &= \vec{g}_i \cdot (\vec{h}_i - \gamma_{i-1} \vec{h}_{i-1}) \\ &= \vec{g}_i \cdot \vec{h}_i - \gamma_{i-1} (\vec{g}_i \cdot \vec{h}_{i-1}).\end{aligned}$$

But $\vec{h}_{i-1} = \vec{g}_{i-1} - \gamma_{i-2} \vec{h}_{i-2}$ so we can get rid of the last term by descent (recall $h_0 = g_0$, which is orthogonal to all the other g_i).

Thus

$$\lambda_i = \frac{\vec{g}_i \cdot \vec{h}_i}{\vec{h}_i^T A \vec{h}_i}, \text{ as claimed.}$$

□

So what? We can construct γ_i without A , but our formula for λ_i still involves A .

Proposition. Suppose $\vec{g}_i = -\nabla f(\vec{P}_i)$ and we proceed in direction \vec{h}_i to the local min at \vec{P}_{i+1} , and let $\vec{g}_{i+1} = -\nabla f(\vec{P}_{i+1})$. Then $\vec{g}_{i+1} = \vec{g}_i - \lambda_i A \vec{h}_i$, following our procedure above.

Proof. We saw last class that if f is quadratic, then $f(\vec{x}) = c - \vec{b} \cdot \vec{x} + \frac{1}{2} \vec{x}^T A \vec{x}$,

$$\nabla f(\vec{x}) = A \vec{x} - \vec{b}.$$

We then have (where A is the Hessian)

$$\vec{g}_i = -\nabla f(\vec{P}_i) = -A \cdot \vec{P}_i + \vec{b}$$

and

$$\begin{aligned}
 (*) \quad g_{i+1} &= -\nabla f(P_{i+1}) \\
 &= -\nabla f(P_i + \lambda \vec{h}_i) \\
 &= -A(P_i + \lambda \vec{h}_i) + b \\
 &= \vec{g}_i - \lambda A \vec{h}_i.
 \end{aligned}$$

where we chose λ to minimize f along the line in direction \vec{h}_i from \vec{P}_i .

But that is exactly the λ which makes

$$\begin{aligned}
 \vec{g}_{i+1} &= -\nabla f(P_{i+1}) \perp \text{ to } h_i, \\
 \text{or}
 \end{aligned}$$

$$g_{i+1} \cdot h_i = 0.$$

or

$$\begin{aligned}
 0 &= \vec{h}_i \cdot (\vec{g}_i - \lambda A \vec{h}_i) \\
 &= \vec{h}_i \cdot \vec{g}_i - \lambda \vec{h}_i A \vec{h}_i
 \end{aligned}$$

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or (wonderfully!)

$$\lambda = \frac{\vec{h}_i \cdot \vec{g}_i}{\vec{h}_i^T A \vec{h}_i} = \lambda_i.$$

□

Thus, just by minimizing in direction h_i and reading off the g_i from ∇f , we can reconstruct the entire sequence of g_i and h_i without ever explicitly knowing A !

We now can outline the conjugate gradient algorithm.

Polak-Ribiere ^{Conjugate Gradient} Algorithm.

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$$\text{Set } p = x_0, \quad g = -\nabla f(p), \quad h = \xi = g.$$

main loop:

minimize from p in direction ξ .

if reduction is small enough, return

$$f_p = f(p), \quad \xi = \nabla f(p).$$

$$\gamma = \frac{(\xi + g) \cdot \xi}{g \cdot g} \quad \left(\gamma = \frac{(g_{i+1} - g_i) \cdot g_{i+1}}{g_i \cdot g_i} \right)$$

$$g = -\xi$$

$$\xi = h = g + \gamma h$$

go back to top of loop.

We note that we could also have

used
$$\gamma = -\frac{\vec{g}_{i+1} \cdot \vec{g}_{i+1}}{\vec{g}_i \cdot \vec{g}_i} \quad (\text{from our corollary})$$

For a function that's exactly quadratic, this would make no difference. But for a non-quadratic function, if we need to do another round of iterations, the Polak-Ribiere version is reported to work better in practice (the other version is called Fletcher-Reeves).

< mathematica demo >